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Protocols for multimedia in the Internet

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Applications and protocol stack

- Which transport protocol?
- UDP is suited for:
 - Request-response (LAN)
 - Multimedia applications
 - Multicast
- TCP (reliability) is suited for:
 - File transfer
 - Terminal emulation
 - Request-response (WAN)
 - Unicast

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UDP protocol

- UDP (User Datagram Protocol) permits application to application (host to host) communication through datagram transmission
- UDP provides a layer 4 service:
 - Connectionless (out of sequence packets)
 - Unreliable (packet lost)
 - Low overhead (slim header)
 - Optional checksum
- Application identification through :
 - Source IP address, destination IP address, source UDP port, destination UDP port
 - No rate control
 - No flow control (possible receiver saturation)
 - No congestion control

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UDP packet format

0	15	16	31
UDP Source Port		UDP Destination Port	
UDP Message Length		UDP Checksum (opt.)	
DATA			

32 bit

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TCP protocol

- TCP (Transmission Control Protocol) is the other Internet layer 4 protocol
- Main characteristics
 - Connection-oriented
 - full-duplex
 - Reliable and in sequence delivery
 - Retransmission
 - Rate control
 - Flow controlled by receiver
 - Congestion control to avoid network saturation

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TCP packet header

0	15	16	31
Source Port Number		Destin. Port Number	
Sequence Number			
Acknowledgement Number			
HLEN	Resv.	flags	Receiver window
Checksum		Urgent Pointer	

32 bit

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Problems for multimedia

- Packetization
 - Voice sample of few bit
 - Single image has very large size
- How to distinguish at the receiver among different coding techniques?
- How to compensate for IP limitations?
 - Packet losses
 - Out of order delivery
 - Packet duplication
- How to notify to the source the correct reception of data?
- How to deal with multicast?

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Using TCP for multimedia?

- TCP is reliable, but
 - Retransmissions cause delays
- TCP is rate controlled to avoid receiver and network congestion, but
 - The available bit rate for the multimedia application is highly variable
- TCP does not support multicast
- TCP cannot be used for real-time multimedia
 - Non real time multimedia can be treated as file transfer

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UDP for multimedia

- UDP supports multicast, is not rate controlled and does not use retransmissions, but
 - Does not guarantee in-order delivery
 - Does not detect and deal with packet losses
 - Does not compensate for delay fluctuations
 - Does not recognize multimedia contents
- IETF proposal: add RTP protocol over UDP

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Example of multimedia application

- IP telephony: three different problems
 - Establish multimedia connection, find IP addresses (possibly multicast), negotiate the type of coding and/or compression scheme, possibly inter-operate with the telephone network
 - Once the connection has been established, transfer audio packets
 - Periodically send feedback information to the transmitter (and to receivers) to indicate the quality of the multimedia connection



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RTP: Real-time Transport Protocol

- Defined in RFC 1889
 - a framework to build multimedia applications
- Provides some basic mechanisms for multimedia data transfer
- It is not an independent protocol, but must be included in the application (no API defined)
- Composed by two different (although related) protocols:
 - RTP: deals with multimedia data transport (even UDP ports)
 - RTCP: provides control and monitoring services (odd UDP ports)
 - Feedback on packet delay and losses
 - Request/response to coding modifications
 - Helps in the management of the list of participants

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RTP + RTCP

- Functions available
 - Payload identification (coding)
 - Sequence number management
 - Timestamp management
 - Monitoring and performance analysis
 - Participants identification
- Functions NOT available
 - No QoS support
 - No reliability, bit rate or delay guarantees
 - No guarantee on correct and in order delivery
 - Exploits UDP checksum to detect errors

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RTP: packet losses

- UDP/IP do not guarantee zero losses
 - Packets may be lost or delivered not in order
- RTP provides sequence numbers in the RTP header to detect out of order delivery

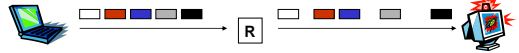
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RTP: delay jitter

- When the multimedia source sends synchronous signals (e.g., voice), one packet is sent every T seconds (source clock)
- The network introduces variable delays even when there are no losses (router buffers)



- How to recover the synchronous signal at the receiver?

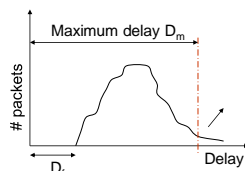
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RTP: delay jitter

- Possible solution: introduce a delay at the receiver
- Use a playback buffer
 - Received packets are stored in the buffer
 - One packet every T seconds is played out
 - The buffer size emulates a fixed delay of D_m seconds
 - D_m is a compromise between low delays and low losses



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RTP: delay jitter

- If the source uses silence suppression?
 - During talk spurt: one packet every T seconds
 - During silence: no packets
- ... a packet may be delayed because
 - the network has delayed it
 - it was preceded by a silence suppressed
- The sequence number is not enough
- A 'timestamp' field (conventional number which increases at the transmitter according to the source clock) is needed in the RTP header to recover information on the generation instant of a packet
 - Timestamp cannot be used as sequence number, since many packets may have the same timestamp (e.g., a video frame generates several packets with the same timestamp)

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RTP: applications

- Applications using RTP may be multicast
- If the network does not support multicast addressing and routing
 - A unicast connection among each pair of participants to the multicast session should be opened
 - The number of connections increases quadratically with respect to the number of users
- An application must specify two RTP parameters:
 - RTP profile
 - A table that associates to each type of payload (coding) a unique code
 - How RTP should use the payload
 - Information such as RTP packet size, number of samples contained in each packet, etc

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RTP: example of voice transmission

- Consider a multi-user voice conference over RTP
- Basic elements:
 - One IP multicast address
 - Users must register the IP multicast address to access the service
 - Two UDP ports
 - One even port for RTP, the next one (odd) for RTCP
 - RTP does not specify how to choose these numbers
- Voice is produced according to a proper coding technique
- Voice samples are grouped in packets
 - Sent to the multicast address
 - The number of samples in each packet depends on the RTP configuration

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RTP: example of voice transmission

- Packets size should be small to keep under control the packetization delay (should be smaller than few tens of ms)
 - Few samples in each packet
- Samples are encapsulated within an IP + UDP + RTP headers



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RTP: example of voice transmission

- Through RTCP, each participants sends (in multicast) statistical data
 - It is possible to analyze service performance
 - Code rate adaptation may be envisioned to adapt the transmission to the measured quality
- Since the RTCP traffic is sent in multicast to all participants (and is generated by all participants) the required bit rate may be significant and should be kept under control

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RTP: example of audio-video conference

- The standard suggests to use two independent RTP flows
 - Advantage: a user may access only one of the two services
- It is mandatory to synchronize the two flows:
 - The RTP timestamps of the two flows may be used with this goal

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RTP: terms and definitions

- An RTP **session**
 - Allows the communication among a set of users through RTP
 - Is identified by a pair of transport addresses
 - A transport address is a IP address-UDP port pair
 - One transport address is used by RTP, the other one for RTCP
 - The IP multicast address (if used) is the same for both pairs
- The **host** or **end-system** (ES) is the user system where applications based on RTP are running. Hosts generate and receive multimedia flows

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RTP: terms and definitions

- An RTP **translator** is a coding translator
 - Modify the data coding and retransmit the modified data flow over the session
 - May also operate as gateways
 - Permit the service provisioning also over non IP islands
- An RTP **mixer** is an RTP flow aggregator
 - More RTP flows in input are mixed and generate a new single RTP flow at output
 - Typically, the bit rate of the new flow is smaller than the sum of the bit rate of original flows
 - Devices more complex than translator
 - Correlated flow synchronization must be dealt with

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RTP: terms and definitions

- RTP **monitors** observe externally the control packet flow (RTCP packets)
 - Do not participate in the RTP transmission/reception process
 - Collect information on the QoS of the RTP session
 - Useful for network providers to control the network service quality

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RTP: terms and definitions

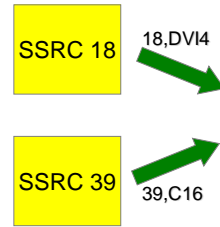
- The **SSRC** (Synchronization SouRCe) is the unique identifier of the data flow generator
 - Is a 32 bit number contained in the RTP header
 - Warning: mixers are characterized by an own SSRC
 - The output data flow from a mixer is a new flow, with its proper timestamp
 - Within an RTP session, each SSRC must be unique
- Since it may be useful or needed to recover the original source of a mixed RTP flow
 - The **CSRC** (Contributing SouRCe) are fields optionally contained in the RTP header that contain the SSRCs of the original sources of the RTP mixed flow

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RTP: Host

- The first host creates an RTP flow with SSRC 18 and DVI4 code, the second host sends RTP data with SSRC 39 and C16 code

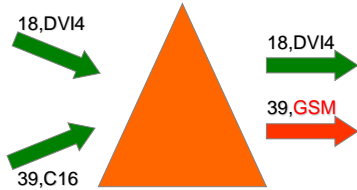


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RTP: Translator

- The translator modifies the coding of the flow with SSRC 39 from C16 to GSM

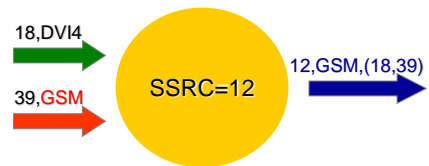


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RTP: Mixer

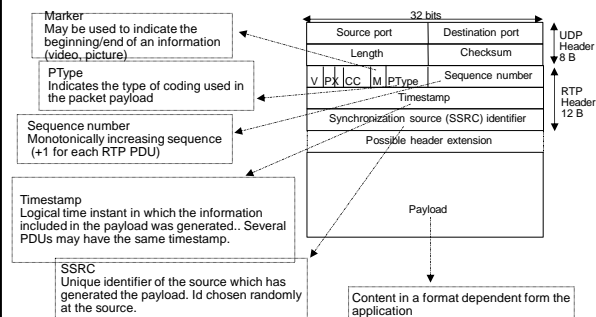
- The RTP flow generated by the mixer has SSRC 12 and CSRC 18 e 39



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RTP packet format



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RTP header

- V (Version) (2bit)
 - Version of the RTP protocol
- P (padding) (1bit)
 - Padding contained in the payload
- X (eXtension header) (1bit)
 - Proprietary
- CC (CSRC count) (4 bit)
 - Number of CSRC field contained in the header
- Marker (1 bit)
 - Use dependent on the RTP session profile
- PT (Payload Type) (7 bit)
 - Type of coding used in the packet payload

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RTP header

- Sequence number (16 bit)
 - The initial sequence number X is chosen randomly at session startup
 - X is inserted in the sequence number field of the first generated RTP packet
 - The second packet will have sequence number X+1, the third X+2...
 - The random extraction minimizes the probability of choosing the same number previously selected in an older RTP session (which may have some packets still in the network)

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RTP header

- Timestamp (32 bit)
 - Represents the "logical" time instant in which the packet payload was generated. More packets may have the same timestamp value
 - The first timestamp value is randomly selected at RTP session startup
 - The timestamp refers to the creation time instant of the first sample contained in the packet
 - Example: if each packet of an RTP phone session contains 160 samples:
 - If packet I has timestamp X, then packet I+1 will have timestamp X+160

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RTP header

- SSRC (32 bit)
 - Identifier of the source that has created the RTP payload. Randomly chosen at the source at session startup. Algorithm defined to solve contentions.
- CSRC (32 bit)
 - Optional field
 - Identifier of the sources that have originally created the payload from which the current payload was derived through a mixing operation

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RTCP: objectives

- Quality of service and congestion control
 - RTCP packets are used as "low frequency" ACKs to signal the reception quality
 - On the basis of RTCP "report", the server may adapt the coding to the communication status
- Identification
 - Provide identification information on RTP session participants (signalling)
- Estimate the number of participants to the multicast session
 - Needed to control the transmission bit rate required by RTCP control signals which would increase too much if the number of participants increases too much

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RTP Control Protocol (RTCP)

- Control protocol for RTP data flow
- Defines the information exchange among the source and the destinations
- Several types of RTCP packets:
 - SR (Sender Report): sent by all active sources to all participants; includes TX and RX statistics
 - RR (Receiver Report): sent by all receivers to all participants; includes RX statistics
 - SDES (Source Description): source description through a unique identifier
 - BYE: session ends or one participant leaves the session
 - APP: application-specific

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RTCP Sender Report

- SR is used to provide information on data recently sent
- A SR contains:
 - An absolute timestamp (NTP timestamp) of the data sending time
 - Relative timestamp referring to the current RT flow
 - Amount of data sent from RTP session start-up
 - Total number of RTP packets sent
 - Total number of byte sent

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RTCP Receiver Report

- RR are sent to inform senders on the quality of the RTP session as seen by receivers
- A RR is sent to each source from which a SR was received
- A RR contains:
 - Identification of the received source
 - Timestamp of the last received SR
 - Delay from the reception of the last SR
 - Highest sequence number received from the source
 - Number of lost RTP packets
 - Fraction of lost RTP packets
 - Estimate of RTP packets jitter

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RTCP SDES

- Used by sources and destinations to identify themselves
- An SDES may contain:
 - CNAME: user identifier (`user@host.domain`)
 - NAME: name of the person using the application
 - EMAIL
 - PHONE
 - LOC: geographical location of the user
 - TOOL: application using RTP
 - NOTES

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RTCP: report transmission speed

- Four rules must be followed to generate RTCP packets
 - The RTCP traffic should be limited to a given percentage of the data traffic (5%)
 - 25% of the session bit rate is devoted to SR packets, the remaining part to other packets
 - An RTCP packet cannot be sent earlier than 5s after the previous RTCP packet transmission
 - A variable time P should be added to the waiting time

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RTCP: report transmission speed

- P is computed as a random number uniformly chosen between 0.5 e 15, multiplied by T_{sr}
- The period for transmitting RTCP packets for the transmitter

$$T_{SR} = \frac{\text{Num_senders}}{0.25 \times 0.05 \times \text{Session_rate}} \times \text{Avr_RTCP_pkt_size}$$

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