

Applications taxonomy

Andrea Bianco
Telecommunication Network Group
firstname.lastname@polito.it
http://www.telematica.polito.it/

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Applications

- Data
 - Generated by single users, by servers, by data centers, by enterprise networks, by P2P architectures, by computing app (e.g. Mapreduce)
 - E-mail, web, messaging, remote login, file transfer, grid computing,
- Voice
- Phone calls, IP calls, skype, ...
- Audio
- Music
- Video
- Multimedia
 - Streaming, videoconferencing

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Applications taxonomy

- · From the bit rate requirements point of view
 - Elastic applications (opportunistic)
 - If resources are available, elastic applications try to exploit them
 - If resources become scarce, elastic applications may reduce their rate (file transfer)
 - Non-elastic applications (multimedia mostly belong to this category)
 - Require a minimum amount of resources
 - If available, the application works properly
 - If not available, the application is unable to work properly
 - May become slightly elastic if changing the coding scheme

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Applications taxonomy

- · From the data loss point of view
 - Loss tolerant
 - · Uncompressed audio, video, voice
 - Loss intolerant
 - File transfer, e-mail, web, grid computing, compressed audio, video, voice
- · From the time sensitivity point of view
 - Not sensitive
 - · File transfer, e-mail, web, grid computing
 - Very sensitive (100ms)
 - Phone
 - Sensitive (few s)
 - Streaming

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Elastic applications

- · Consider a file transfer
- Small end-to-end delay preferred (not fundamental)
- · Required bit rate: the higher the better but it may be low
- Packet losses recovered by the transport protocol through retransmission (less often through error correction)
 - End-to-end delay increases

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Real-time multimedia and streaming

- Real time multimedia applications
- Two users interact (in real time)
- Low delay fundamental (a delayed packet is equivalent to a lost packet)
- Required bit-rate may be significant depending on whether video is involved
- May be robust to (limited) packet losses depending on the compression level
- Multimedia streaming applications
- No real time requirements
- May tolerate packet delays if initial delay large (buffering)
- Required bit-rate may be significant depending on whether video is involved or not
- May be robust to (limited) packet losses depending on the compression level

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Example of real time multimedia application

IP telephony: three different problems

Establish multimedia connection, find IP addresses (possibly multicast), negotiate the type of coding and/or compression scheme, possibly inter-operate with the telephone network

Once the connection has been established, transfer audio packets

RTCP

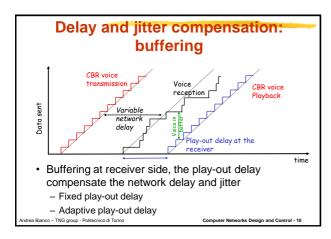
Periodically send feedback information to the transmitter (and to receivers) to indicate the quality of the (possibly multimedia) connection

Real-time multimedia: Internet Phone Voice as input: sounds and silence period alternate

- Packets generated at a constant rate or when the source emitting power is above a given threshold:
 - E.g.: 20 ms of voice sample at 8kb/s
- Packets are delayed (should be compensated) and lost:
 - Network losses, due to congestion
 - · max tolerable may be 10%
 - Losses due to excessive delays (IP datagram received too late for
 - Max tolerable is roughly 400 ms
- · Compensation techniques
 - At the transmitter (adaptive coding)
 - At the receiver (buffering)

Reaction to losses, delay and jitter

- · Use of a variable bit-rate coder
 - Send small size packets when congestion is detected and the experienced delay is high
 - Send large size packets if the network is lightly loaded
- Quality of reception estimate mechanisms are needed
- The transmitter bit rate should be controlled according to:
 - Instantaneous and/or average loss rate
 - Absolute or relative delay
 - Delay jitter

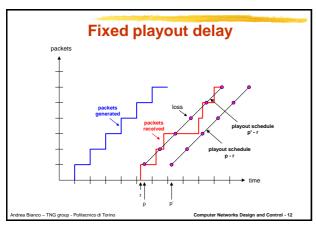


Fixed playout delay

- The receiver plays out each voice sample exactly q seconds after the sample generation
 - If the sample has timestamp t, it is played out at t+q
 - If the sample is received after t+q, it is considered as lost
 - Coding scheme may compensate for losses
- The value assumed by q:
 - if q is large: less packets are lost, higher delay, more buffering needed
 - if q is small: improved interactivity

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Adaptive playout delay

- Objective: minimize play-out delay while keeping low the loss rate
 - Estimate the network delay, to determine the play-out delay at speech startup
 - Compress or extend the silence periods
 - Samples always reproduced after 20ms during activity periods

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Multimedia streaming

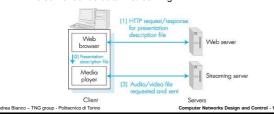
- Streaming
 - Multimedia file stored at the source
 - Sent to the client
 - File play-out starts when the file transfer is under way
 - Constraint: missing data should reach the receiver before the play-out ends
 - Alternative to file download to playback it later (file transfer!)

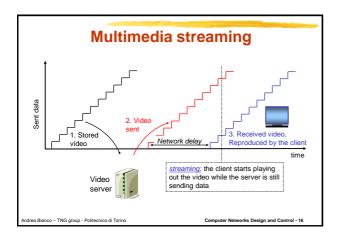
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Multimedia: streaming approach

- The client browser receives the metafile containing the multimedia streaming file description
- · The browser passes the metafile to the player
- · The player contacts the streaming server
- · The server sends data in streaming





Multimedia streaming with buffering at the client side • Tradeoff betweeen initial delay (buffer size) and tolerance to network jitter client buffer drain rate = d from network fill rate = x(t) network Andrea Blanco - TNG group - Politecnico di Torino Computer Networks Design and Control - 17