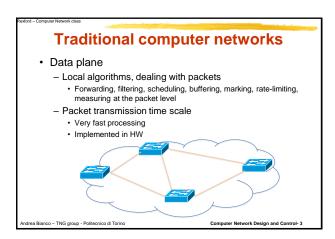
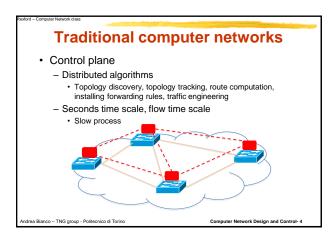
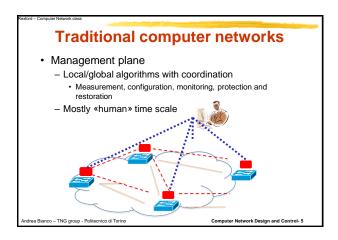


Outline

- SDN
 - Motivations and definitions
 - Centralized architecture
 - Flow based forwarding
- Openflow protocol
- Advances
 - Distributed controllers
 - Stateful switches

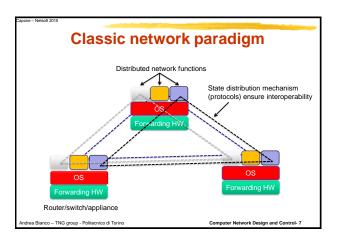


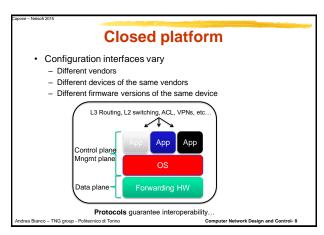


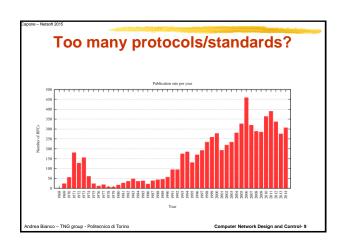


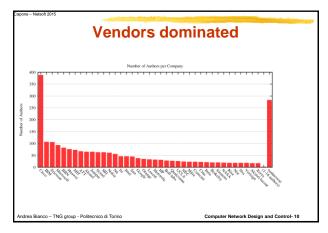
Traditional computer networks

- - Incredible success (from research experiments to global commercial infrastructure)
 - «In principle» complexity at the edge
 - · «Only» packet forwarding inside
 - · Complexity at the edge (SW) enables fast innovation
 - Host running increasingly complex applications (SW)
 Web, P2P, social networks, virtual reality, video streaming
 - Inside the network?
 - · Closed equipments, SW and HW intermixed, vendor specific interfaces, many more features beside forwarding, too many
 - · Slow and costly development and management





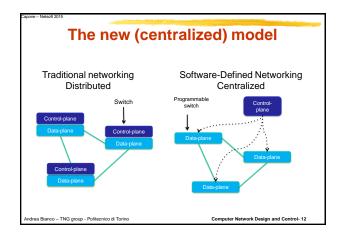


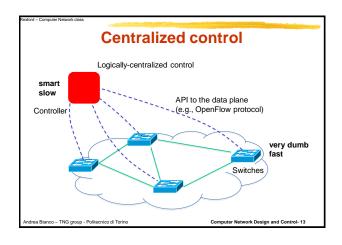


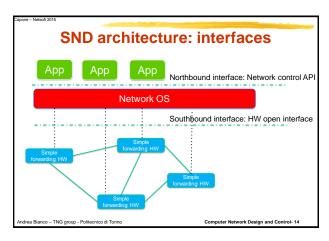
Software Defined Networking*

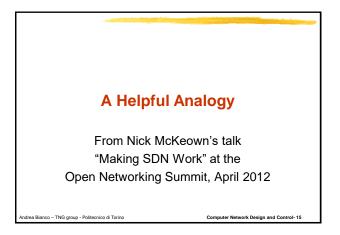
- · "New" key elements
 - Clean interface (API) between data and control plane
 - Logically centralized control plane
 - Control plane out of forwarding devices
 - Control plane (SW) may run on general purpose HW
 - Global network view
 - SDN controller or Network Operating Systems
 - Network programmability
 - New architecture
 - Flow based switching
 - · Programmed by the centralized controller
 - · Very flexible flow definition
 - Network applications running on top of NOS

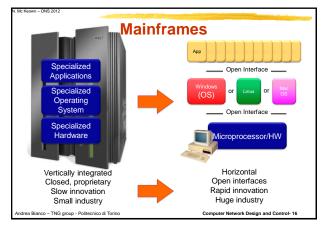
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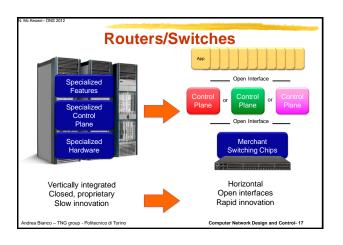


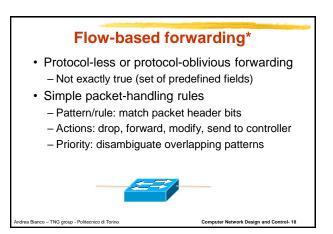


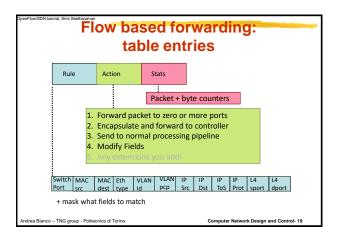


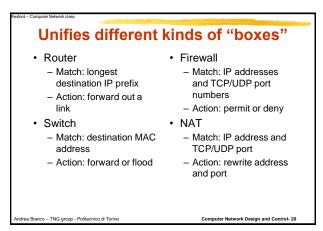


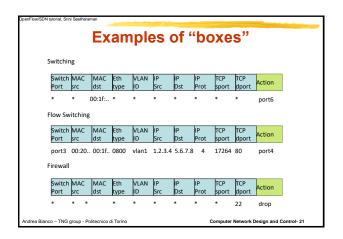


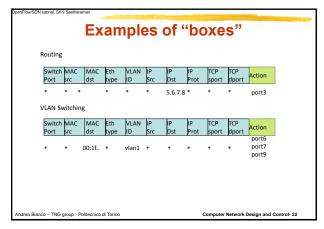


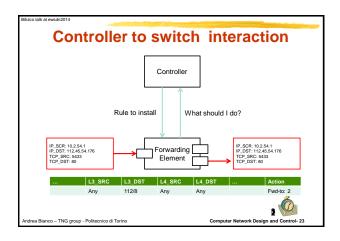


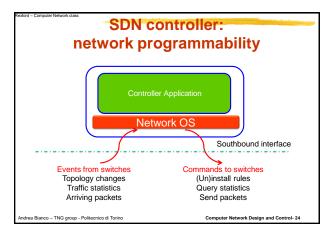












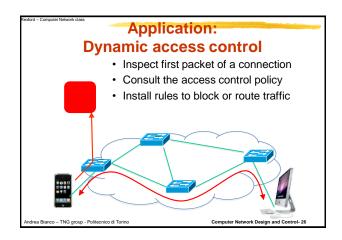
Example of applications

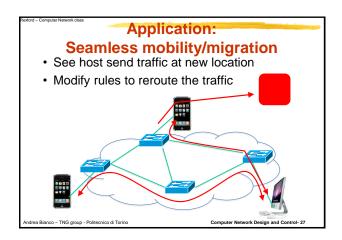
- · Dynamic access control
- · Seamless mobility/migration
- · Server load balancing
- · Network virtualization
- · Using multiple wireless access points
- · Traffic engineering
- · Energy-efficient networking
- · Adaptive traffic monitoring
- · Denial-of-Service attack detection

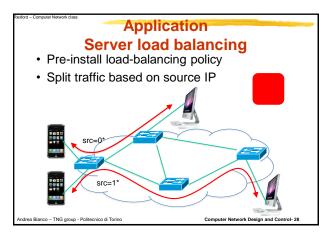
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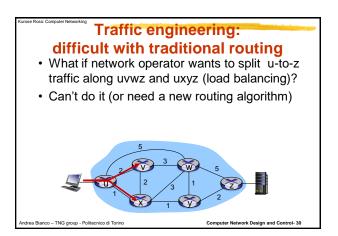
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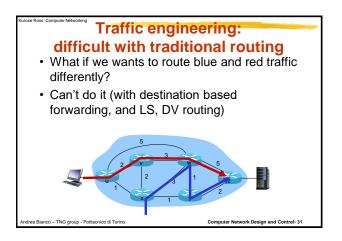


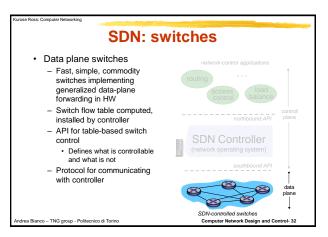


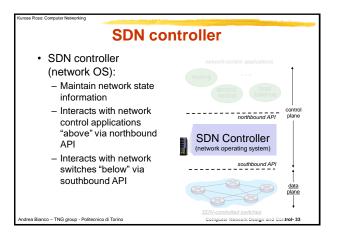


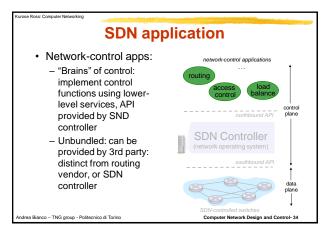
Traffic engineering: difficult with traditional routing Hp. Destination based routing • What if network operator wants - u-to-z traffic to flow along uvwz - x-to-z traffic to flow xwyz? • Need to define link weights so traffic routing algorithm computes routes (or need a new routing algorithm) • Does not work - Modifies many routes - Cannot change weights to route each individual flow

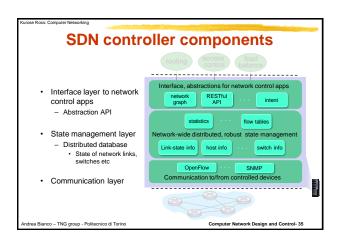


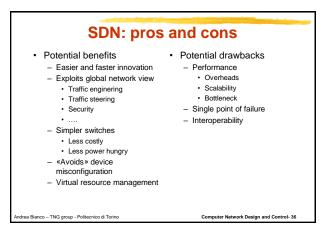












SDN where?

- · Campus LAN
- · Data center
- · WAN (google) to interconnect data centers
- ISP?
- 5G networks

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The role of the scenario

- Datacenter
 - Very large number of devices
 - · Spatially collocated
 - Low and predictable delays between devices
 - Dedicated network for control
 - Out of band control traffic
- · ISP/POP
 - Lower number of devices
 - · Spatially distributed
 - High and unpredictable latencies
 - Control and data share the same resources
 - · In band control traffic

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Level of aggregation

- · Flow Based
- Group Based
- Every flow is individually set up by controller
- Exact-match flow entries
- Flow table contains one entry per flow
- Suited for fine grain control, e.g. campus networks
- One flow entry covers
- large groups of flows

 Wildcard flow entries
- Flow table contains one entry per category/group of flows
- Suited for large number of flows, e.g. ISPs

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Level of aggregation

- · High aggregation level
 - Dealing with few large objects
 - Reduced occupation of forwarding table
 - Reduced signaling overhead and controller load
 - Coarse granularity in the control of flow Qos
 - A flow steering moves a large amount of traffic
 - Less elements to deal with for load balancing but more difficult to balance

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Reactive vs. Proactive

- Reactive
 - Flow table empty at boot
 - First packet of a flow sent to the controller
 - Controller inserts flow entries
 - Dynamic network
 - Every flow incurs small (?) additional flow setup time
 - Large control traffic
 - Large load on the controller
 Efficient use of flow table
 - If control connection lost, switch has limited utility

- Proactive
 - Controller pre-populates flow table in switch at boot
 - Zero additional flow setup time
 - Static network
 - Loss of control connection does not disrupt traffic
 - Essentially requires
 - aggregated (wildcard) rules

 Reduced table size

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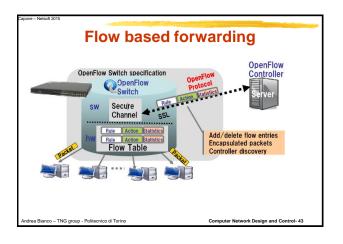
OpenFlow protocol

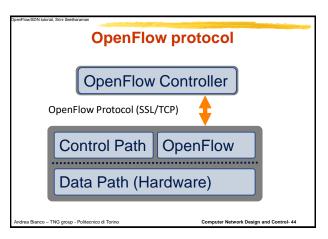
Andrea Bianco andrea.bianco@polito.it http://www.telematica.polito.it/

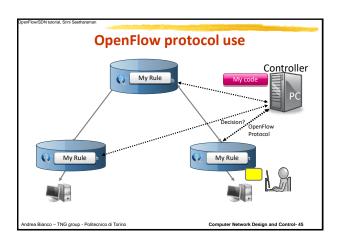
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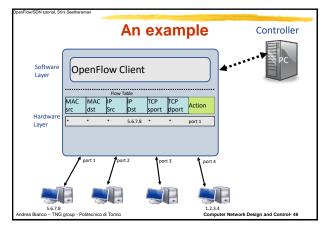
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OpenFlow protocol messages

- · Controller-to-switch
 - Initiated by the controller and used to directly manage or inspect the state of the switch
 - Features, Config, Modify State, Read State, Packet Out, Barrier
- · Asynchronous
 - Sent to the controller without controller soliciting
 - Packet-in, Flow Removed/Expiration, Port status, Error, ...
- Symmetric
 - Sent without solicitation in any direction
 - Hello, Echo, Experimenter/Vendor

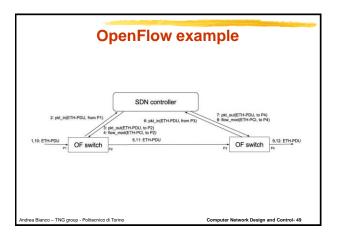
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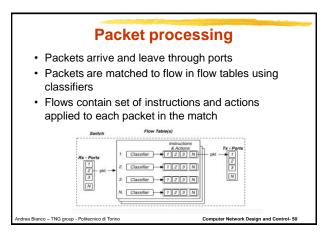
OpenFlow (main) messages

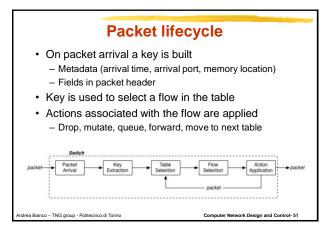
- Packet_in
 - Switch to controller
 - Carries a packet copy (possibly only the header)
 - · What is best?
 - Generated by default in case of table miss
- Packet_out
 - Controller to switch
 - Send the packet out of a specified port
 - Carries the full packet or the switch buffer id
- Flow_mod
 - Controller to switch
 - Modify flow tables
 - Carries match-action rule to install

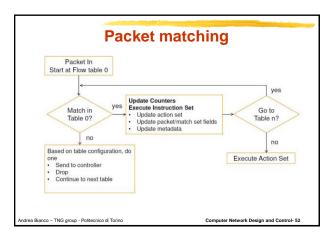
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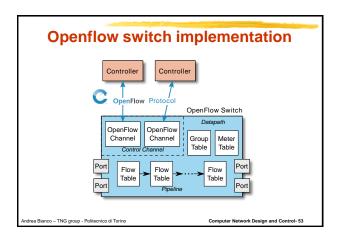
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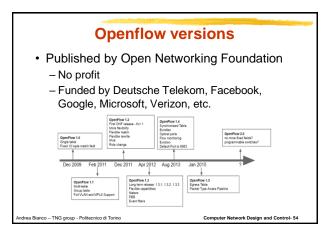




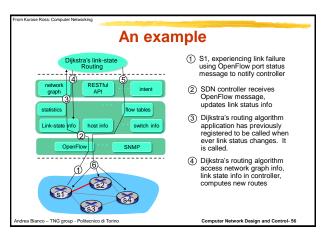


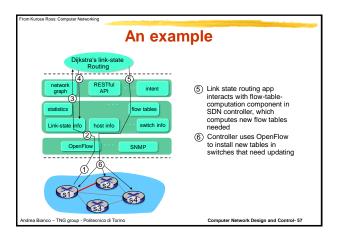




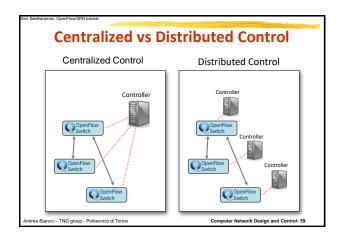












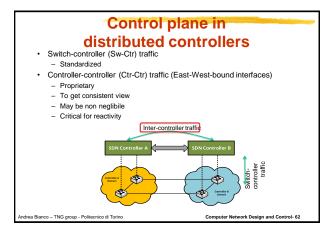
Why distributed/multiple controllers? • To enhance resilience to failures - Controller failures can be managed - Still to deal with failures in data and control plane • To solve scalability issues - Faster controllers • Limited scaling - More proactive rules to reduce number of requests • Limited flexibility - Multiple controllers • Permit load balancing to reduce processing load • Permit switch migration

Distributed controllers

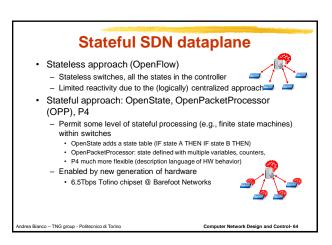
- · Virtual topology among controllers
 - to coordinate the operations of the controllers
 - peer, hierarchical, master/slave
- · Network view maintenance
 - different levels of consistency (strong/weak) among the controllers
 - affects the reactivity
 - may lead to temporary rule conflicts

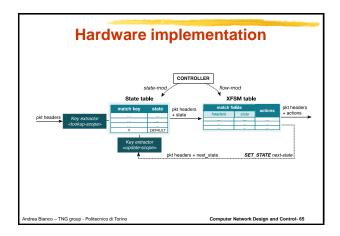
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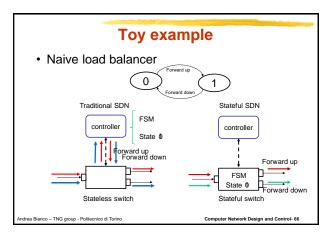
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Stateful data plane Andrea Bianco andrea.bianco@polito.it http://www.telematica.polito.it/







Traffic classification Mirror a pre-defined number of packets to traffic classifier for each flow Interrupt the mirroring if the flow is identified Taffic course description of the flow is identified. Taffic course description of the flow is identified.

Stateful benefits

- · Improve network reactivity
 - Simple local decisions at the switch
 - Reduced controller load
 - Reduced signaling overhead
- Permits to gracefully move functionalities
 - Balance central vs distributed control
- · Not all switches need to be stateful
 - State positioning or distribution

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