



## Software Defined Networking (SDN)

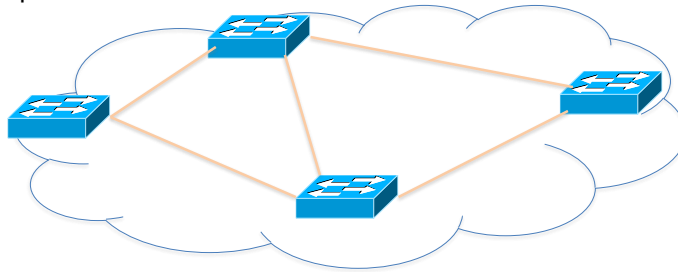
Andrea Bianco  
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<http://www.telematica.polito.it/>

## Outline

- SDN
  - Motivations and definitions
  - Centralized architecture
  - Flow based forwarding
- Openflow protocol
- Advances
  - Distributed controllers
  - Stateful switches

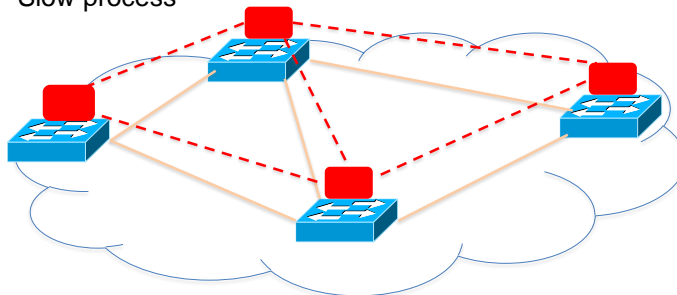
## Traditional computer networks

- Data plane
  - Local algorithms, dealing with packets
    - Forwarding, filtering, scheduling, buffering, marking, rate-limiting, measuring at the packet level
  - Packet transmission time scale
    - Very fast processing
    - Implemented in HW



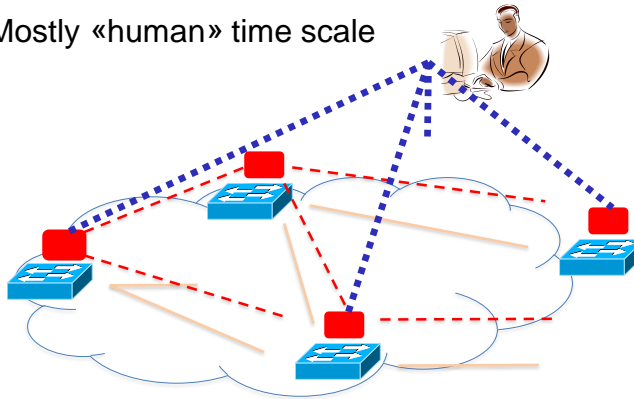
## Traditional computer networks

- Control plane
  - Distributed algorithms
    - Topology discovery, topology tracking, route computation, installing forwarding rules, traffic engineering
  - Seconds time scale, flow time scale
    - Slow process



## Traditional computer networks

- Management plane
  - Local/global algorithms with coordination
    - Measurement, configuration, monitoring, protection and restoration
  - Mostly «human» time scale

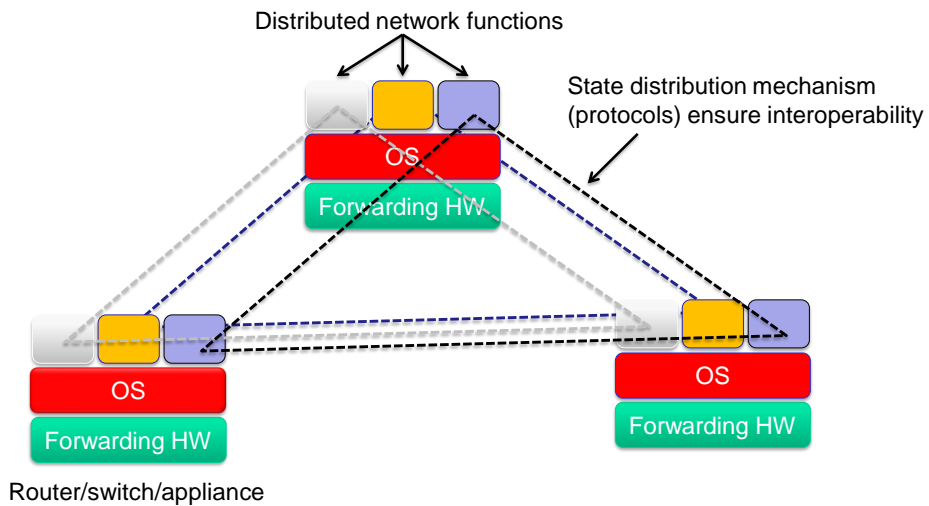


## Traditional computer networks

- Features
  - Incredible success (from research experiments to global commercial infrastructure)
  - «In principle» complexity at the edge
    - «Only» packet forwarding inside
    - Complexity at the edge (SW) enables fast innovation
    - Host running increasingly complex applications (SW)
      - Web, P2P, social networks, virtual reality, video streaming
  - Inside the network?
    - Closed equipments, SW and HW intermixed, vendor specific interfaces, many more features beside forwarding, too many protocols
    - Slow and costly development and management

Capone – Netsoft 2015

## Classic network paradigm



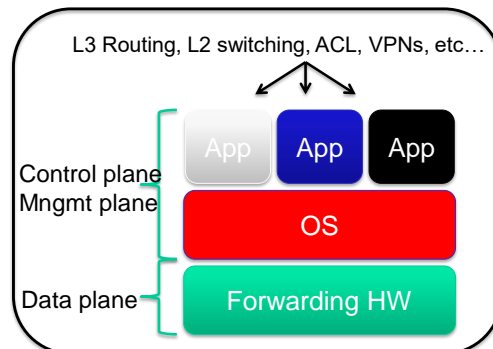
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## Closed platform

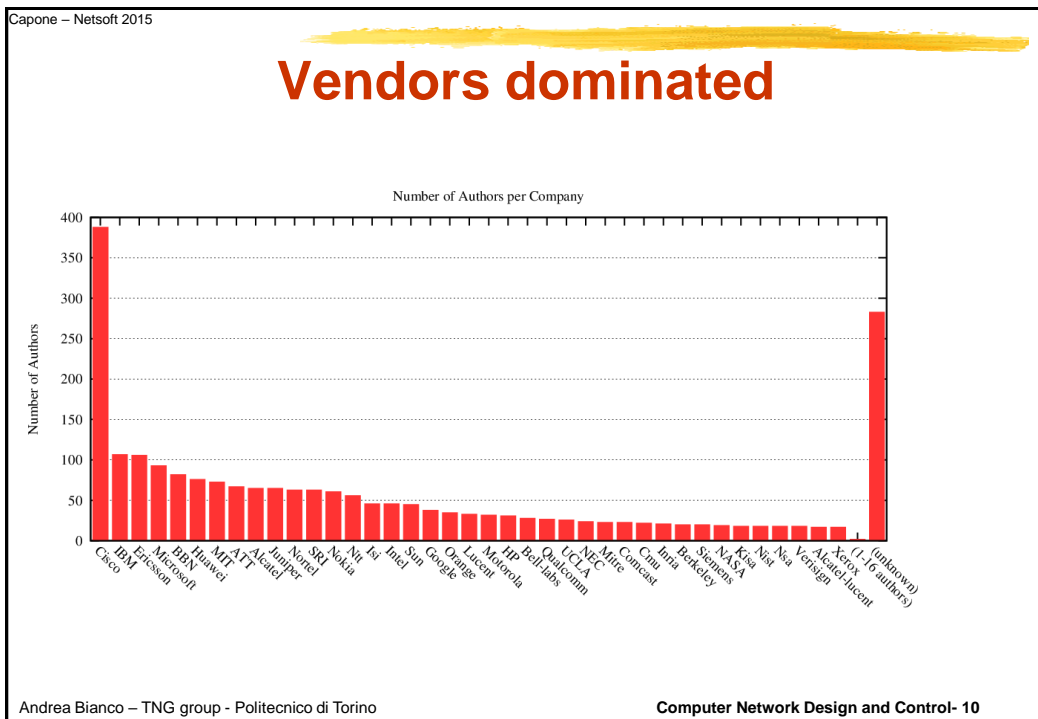
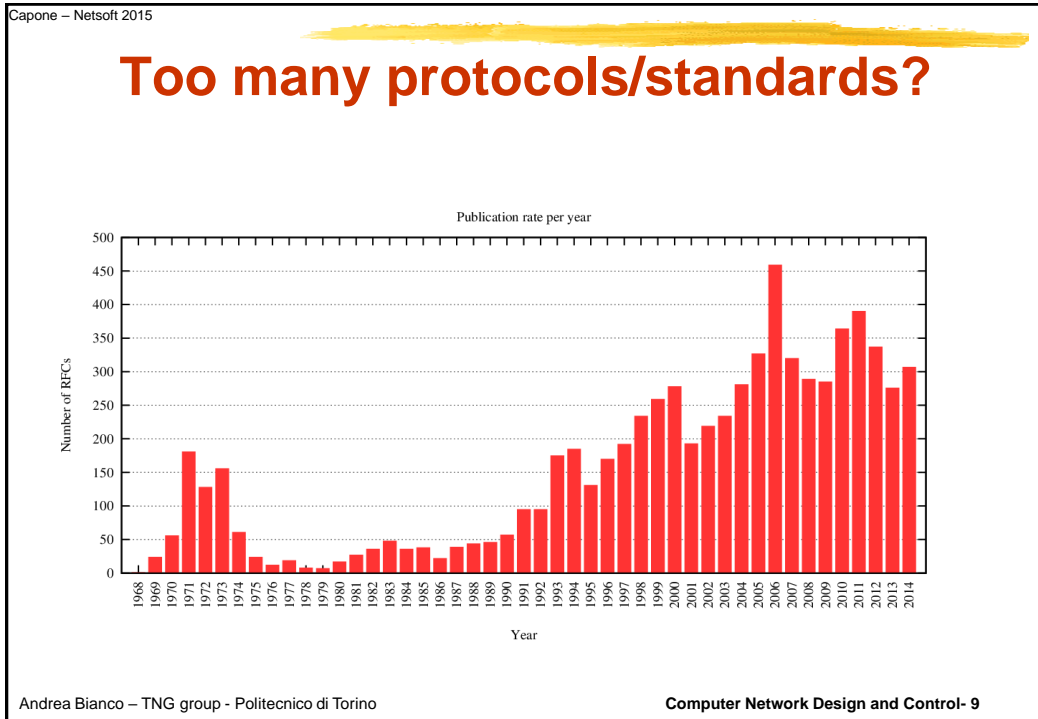
- Configuration interfaces vary
  - Different vendors
  - Different devices of the same vendors
  - Different firmware versions of the same device



Protocols guarantee interoperability...

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## Software Defined Networking\*

- “New” key elements
  - Clean interface (API) between data and control plane
  - Logically centralized control plane
    - Control plane out of forwarding devices
    - Control plane (SW) may run on general purpose HW
    - Global network view
    - SDN controller or Network Operating Systems
      - Network programmability
      - New architecture
  - Flow based switching
    - Programmed by the centralized controller
    - Very flexible flow definition
  - Network applications running on top of NOS

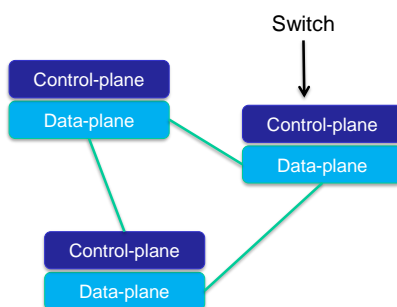
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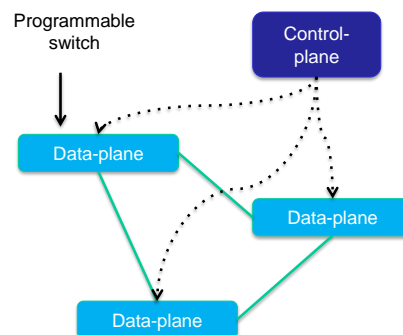
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## The new (centralized) model

Traditional networking  
Distributed

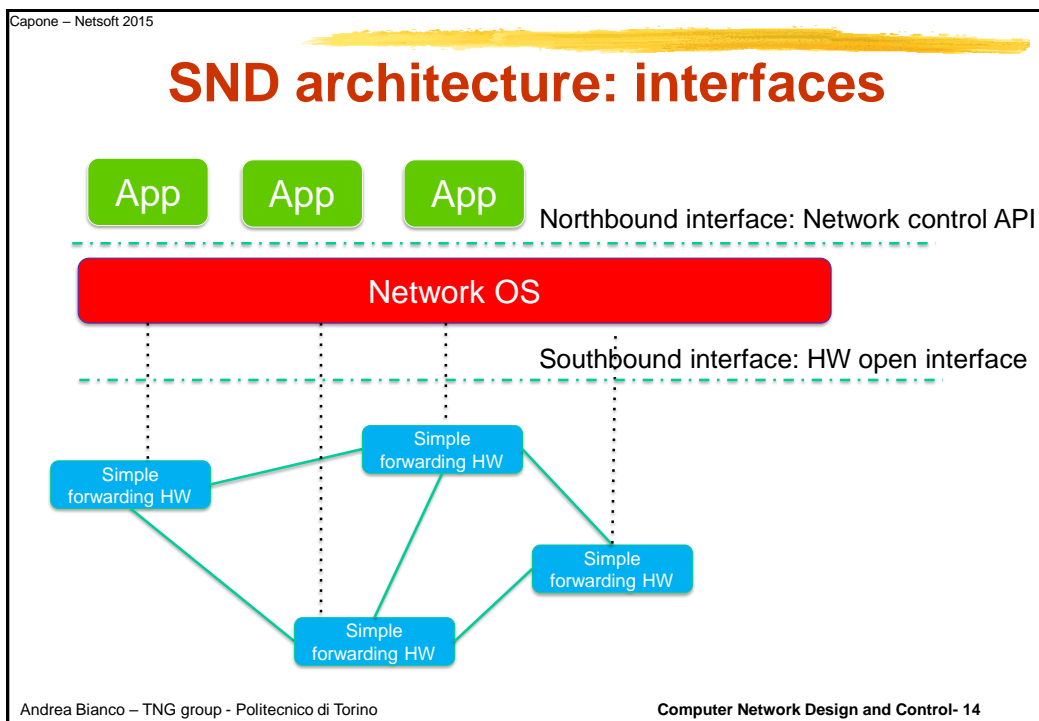
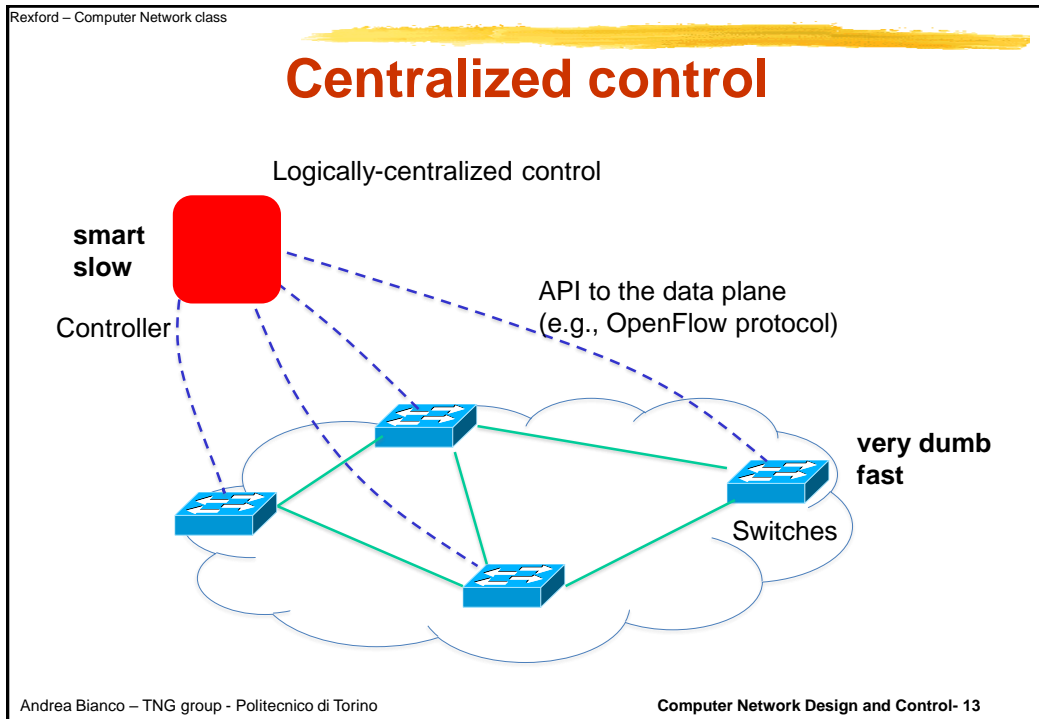


Software-Defined Networking  
Centralized



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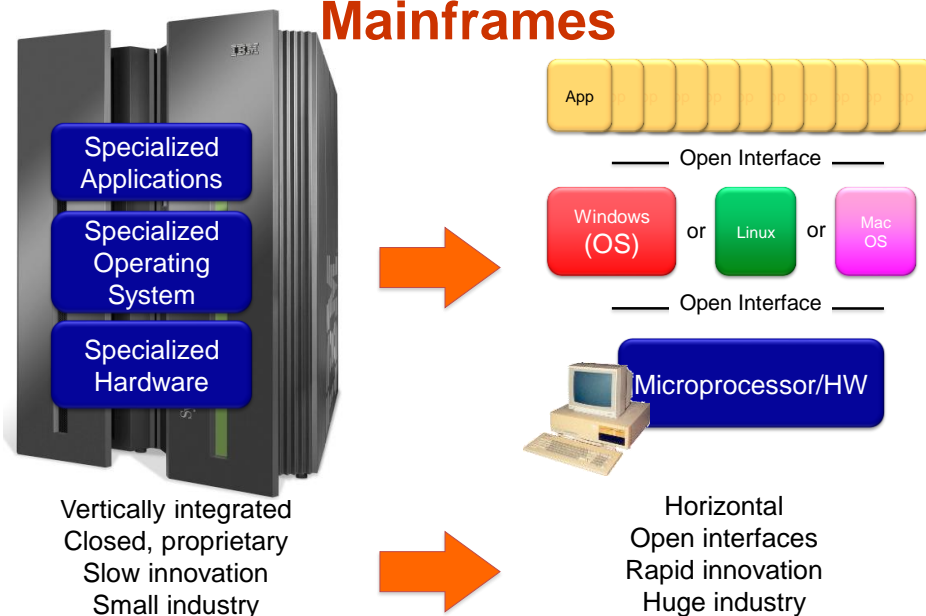
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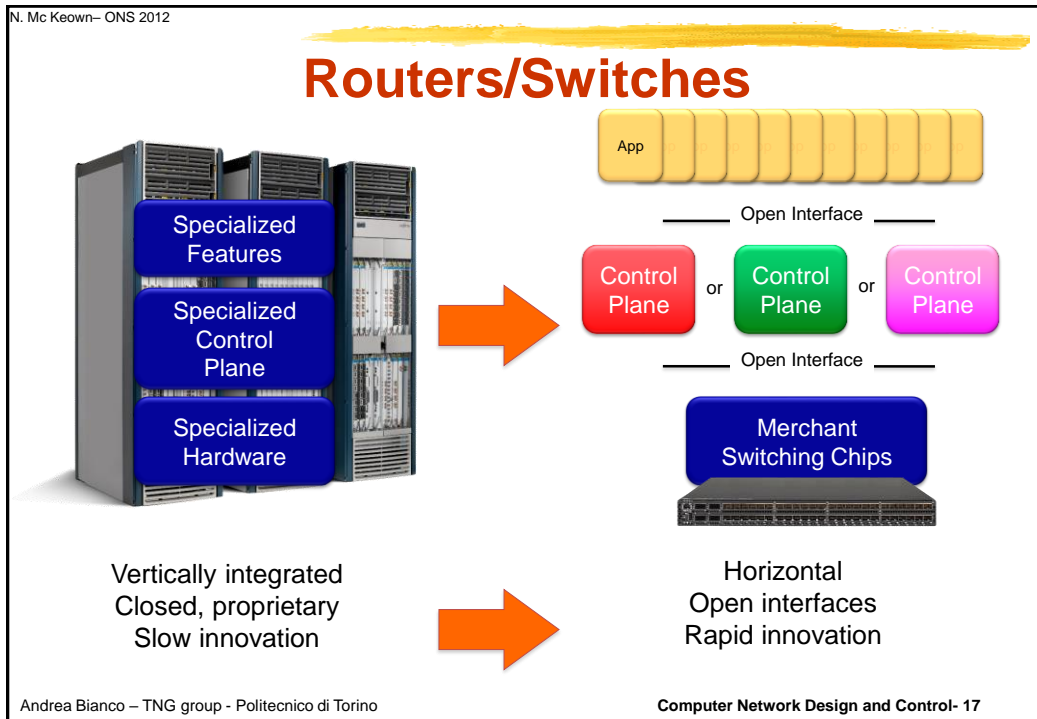
## A Helpful Analogy

From Nick McKeown's talk  
"Making SDN Work" at the  
Open Networking Summit, April 2012

## Mainframes







## Flow-based forwarding\*

- Protocol-less or protocol-oblivious forwarding
  - Not exactly true (set of predefined fields)
- Simple packet-handling rules
  - Pattern/rule: match packet header bits
  - Actions: drop, forward, modify, send to controller
  - Priority: disambiguate overlapping patterns



OpenFlow/SDN tutorial, Srin Seetharaman

## Flow based forwarding: table entries

Rule	Action	Stats
------	--------	-------

Packet + byte counters

1. Forward packet to zero or more ports
2. Encapsulate and forward to controller
3. Send to normal processing pipeline
4. Modify Fields
5. Any extensions you add!

Switch Port	MAC src	MAC dest	Eth type	VLAN Id	VLAN pcp	IP Src	IP Dst	IP ToS	IP Prot	L4 sport	L4 dport
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+ mask what fields to match

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## Unifies different kinds of “boxes”

- Router
  - Match: longest destination IP prefix
  - Action: forward out a link
- Switch
  - Match: destination MAC address
  - Action: forward or flood
- Firewall
  - Match: IP addresses and TCP/UDP port numbers
  - Action: permit or deny
- NAT
  - Match: IP address and TCP/UDP port
  - Action: rewrite address and port

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## Examples of "boxes"

### Switching

Switch Port	MAC src	MAC dst	Eth type	VLAN ID	IP Src	IP Dst	IP Prot	TCP sport	TCP dport	Action
*	*	00:1f:..	*	*	*	*	*	*	*	port6

### Flow Switching

Switch Port	MAC src	MAC dst	Eth type	VLAN ID	IP Src	IP Dst	IP Prot	TCP sport	TCP dport	Action
port3	00:20:..	00:1f:..	0800	vlan1	1.2.3.4	5.6.7.8	4	17264	80	port4

### Firewall

Switch Port	MAC src	MAC dst	Eth type	VLAN ID	IP Src	IP Dst	IP Prot	TCP sport	TCP dport	Action
*	*	*	*	*	*	*	*	*	22	drop

## Examples of "boxes"

### Routing

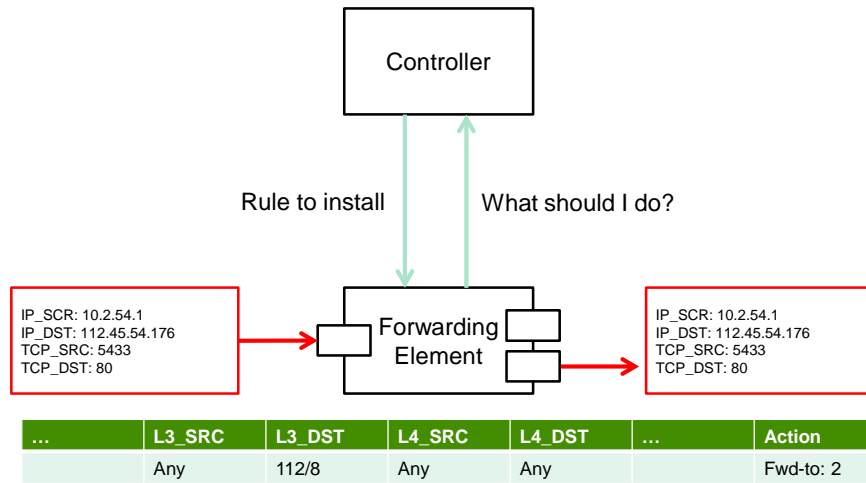
Switch Port	MAC src	MAC dst	Eth type	VLAN ID	IP Src	IP Dst	IP Prot	TCP sport	TCP dport	Action
*	*	*	*	*	*	5.6.7.8	*	*	*	port3

### VLAN Switching

Switch Port	MAC src	MAC dst	Eth type	VLAN ID	IP Src	IP Dst	IP Prot	TCP sport	TCP dport	Action
*	*	00:1f:..	*	vlan1	*	*	*	*	*	port6 port7 port9

Bifulco talk at ewsdn2014

## Controller to switch interaction

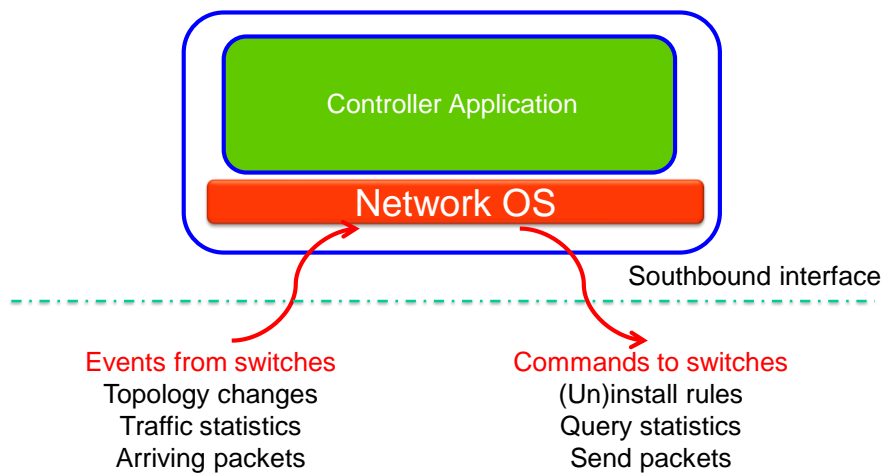


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## SDN controller: network programmability



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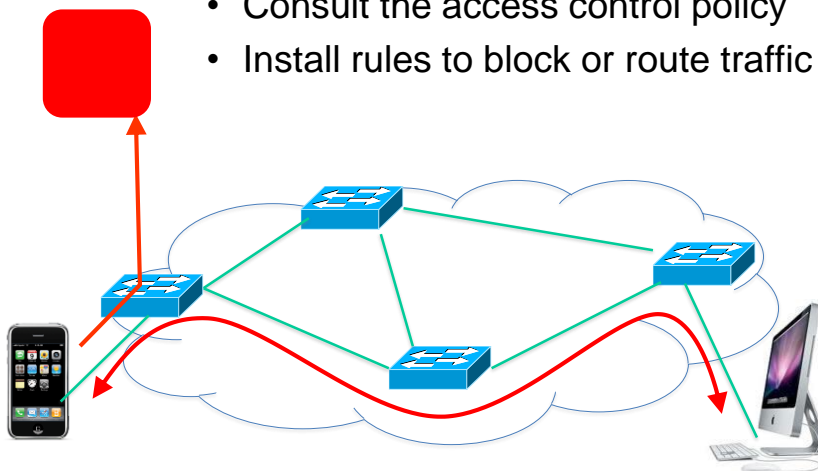
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## Example of applications

- Dynamic access control
- Seamless mobility/migration
- Server load balancing
- Network virtualization
- Using multiple wireless access points
- Traffic engineering
- Energy-efficient networking
- Adaptive traffic monitoring
- Denial-of-Service attack detection
- .....

## Application: Dynamic access control

- Inspect first packet of a connection
- Consult the access control policy
- Install rules to block or route traffic

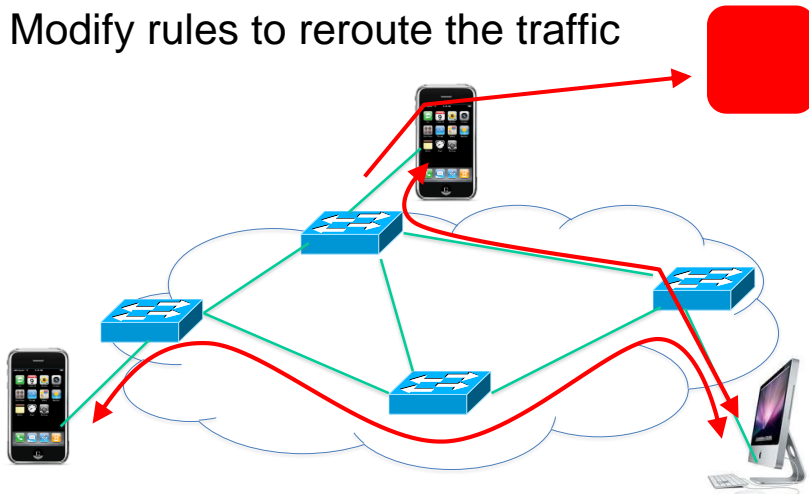


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## Application:

### Seamless mobility/migration

- See host send traffic at new location
- Modify rules to reroute the traffic



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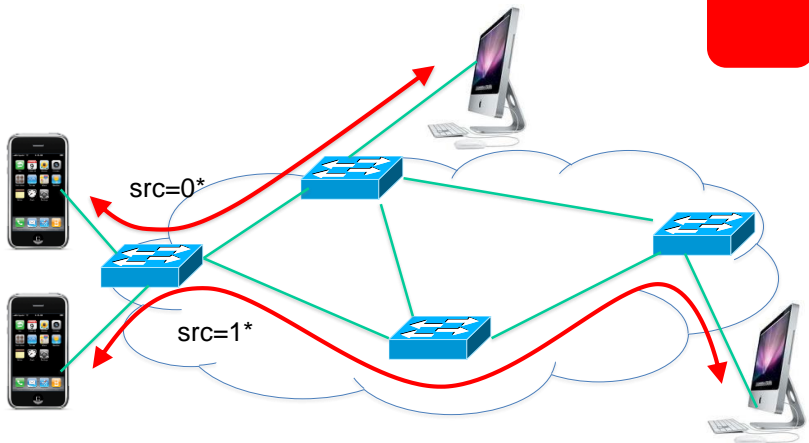
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## Application

### Server load balancing

- Pre-install load-balancing policy
- Split traffic based on source IP



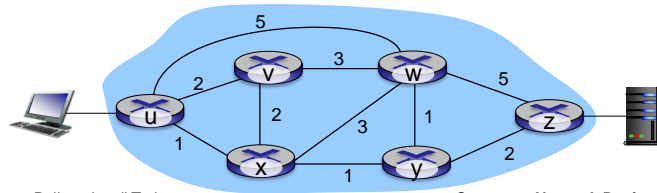
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## Traffic engineering: difficult with traditional routing

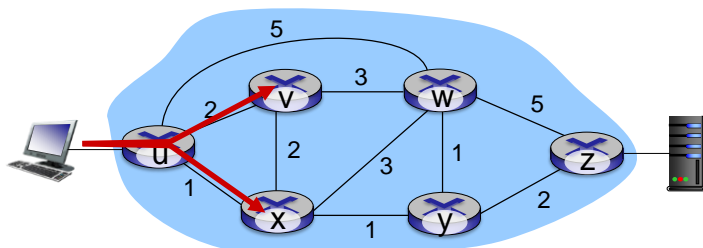
Hp. Destination based routing

- What if network operator wants
  - u-to-z traffic to flow along uvwz
  - x-to-z traffic to flow xyz?
- Need to define link weights so traffic routing algorithm computes routes (or need a new routing algorithm)
- Does not work
  - Modifies many routes
  - Cannot change weights to route each individual flow



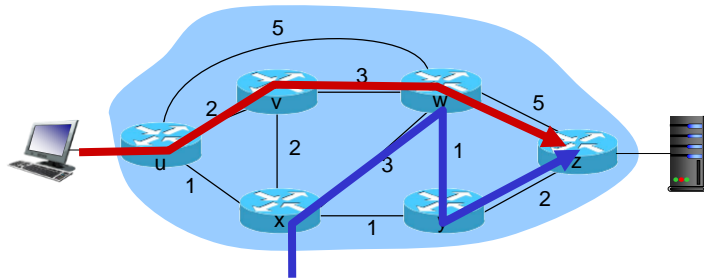
## Traffic engineering: difficult with traditional routing

- What if network operator wants to split u-to-z traffic along uvwz and uxyz (load balancing)?
- Can't do it (or need a new routing algorithm)



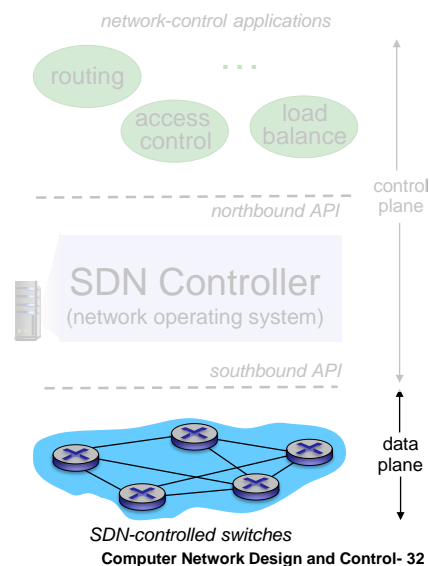
## Traffic engineering: difficult with traditional routing

- What if we want to route blue and red traffic differently?
- Can't do it (with destination based forwarding, and LS, DV routing)



## SDN: switches

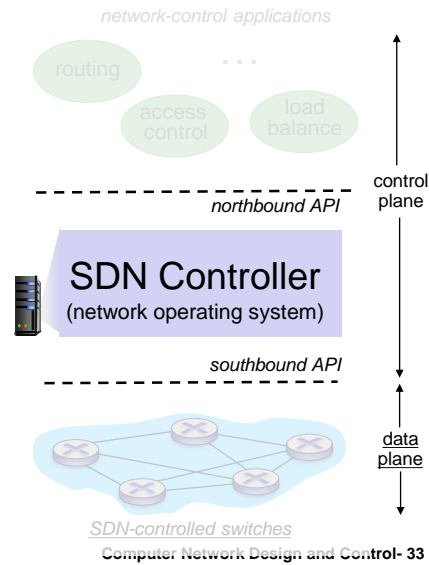
- Data plane switches
  - Fast, simple, commodity switches implementing generalized data-plane forwarding in HW
  - Switch flow table computed, installed by controller
  - API for table-based switch control
    - Defines what is controllable and what is not
  - Protocol for communicating with controller





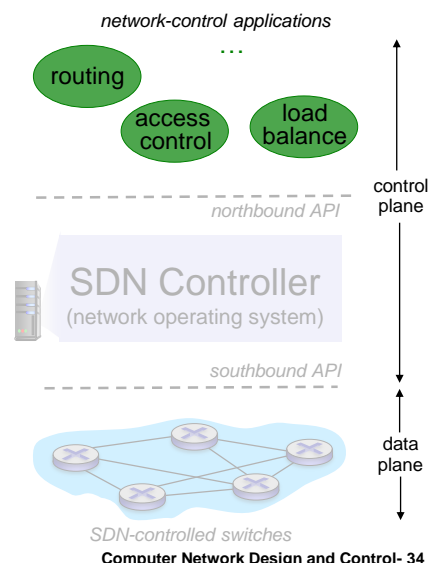
## SDN controller

- SDN controller (network OS):
  - Maintain network state information
  - Interacts with network control applications “above” via northbound API
  - Interacts with network switches “below” via southbound API



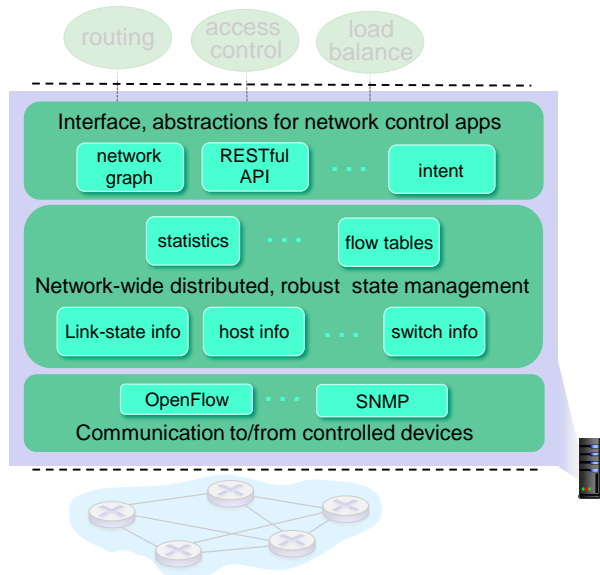
## SDN application

- Network-control apps:
  - “Brains” of control: implement control functions using lower-level services, API provided by SDN controller
  - Unbundled: can be provided by 3rd party: distinct from routing vendor, or SDN controller



## SDN controller components

- Interface layer to network control apps
  - Abstraction API
- State management layer
  - Distributed database
    - State of network links, switches etc
- Communication layer



## SDN: pros and cons

- Potential benefits
  - Easier and faster innovation
  - Exploits global network view
    - Traffic engineering
    - Traffic steering
    - Security
    - ....
  - Simpler switches
    - Less costly
    - Less power hungry
  - «Avoids» device misconfiguration
  - Virtual resource management
- Potential drawbacks
  - Performance
    - Overheads
    - Scalability
    - Bottleneck
  - Single point of failure
  - Interoperability

## SDN where?

- Campus LAN
- Data center
- WAN (google) to interconnect data centers
- ISP?
- 5G networks

## The role of the scenario

- Datacenter
  - Very large number of devices
    - Spatially collocated
  - Low and predictable delays between devices
  - Dedicated network for control
    - Out of band control traffic
- ISP/POP
  - Lower number of devices
    - Spatially distributed
  - High and unpredictable latencies
  - Control and data share the same resources
    - In band control traffic



## Level of aggregation

- Flow Based
  - Every flow is individually set up by controller
  - Exact-match flow entries
  - Flow table contains one entry per flow
  - Suited for fine grain control, e.g. campus networks
- Group Based
  - One flow entry covers large groups of flows
  - Wildcard flow entries
  - Flow table contains one entry per category/group of flows
  - Suited for large number of flows, e.g. ISPs

## Level of aggregation

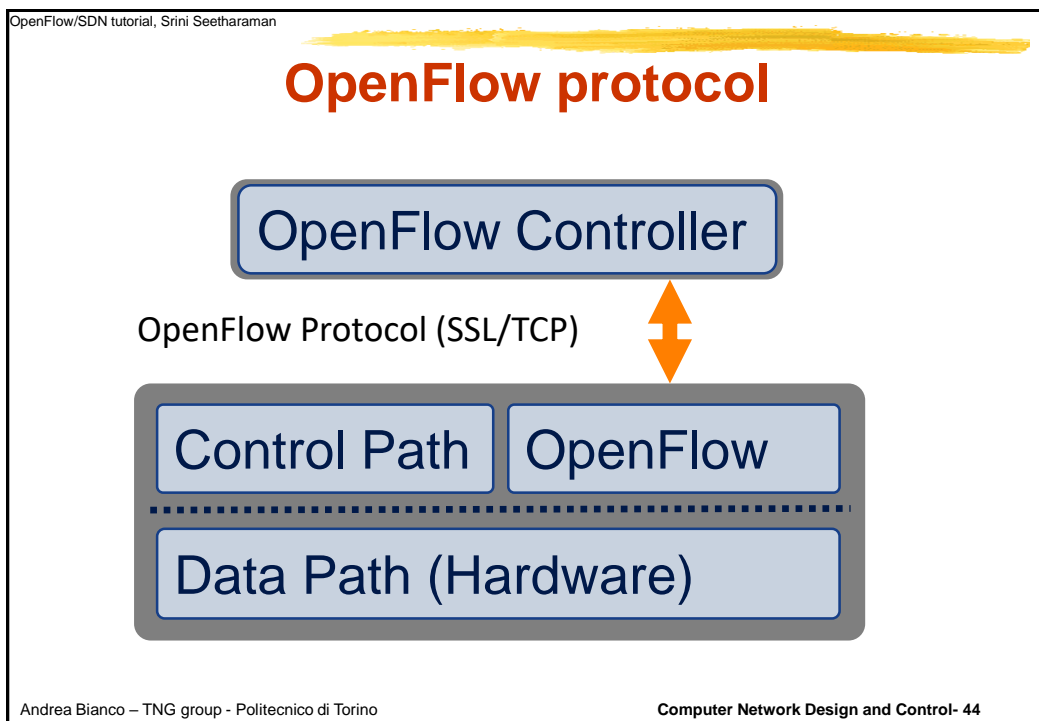
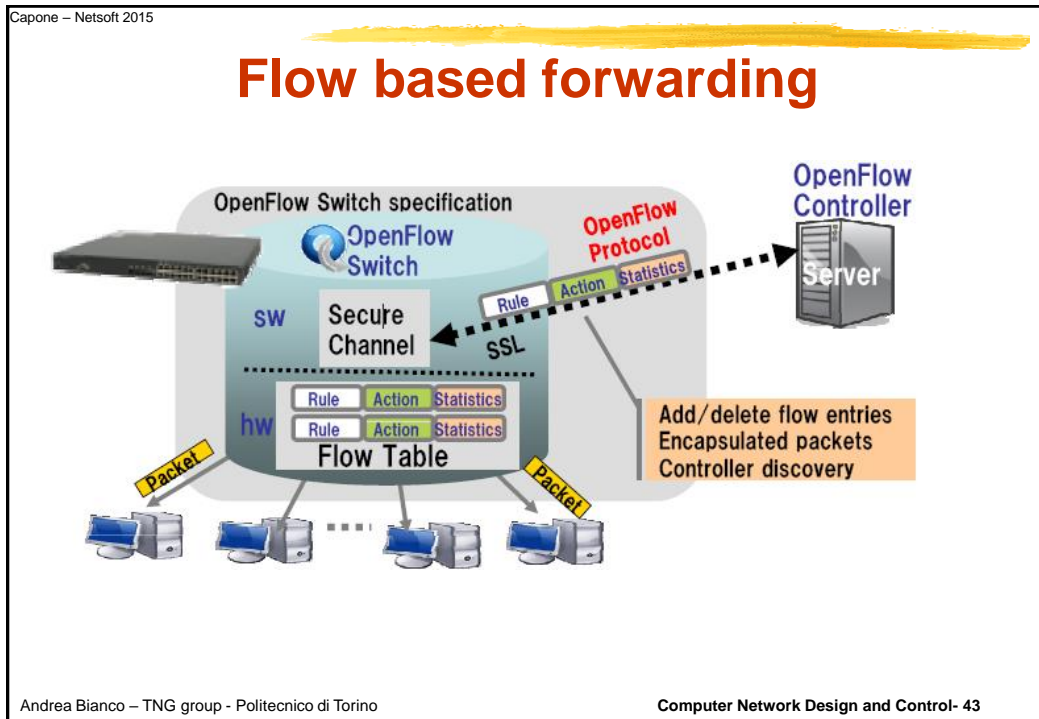
- High aggregation level
  - Dealing with few large objects
  - Reduced occupation of forwarding table
  - Reduced signaling overhead and controller load
  - Coarse granularity in the control of flow Qos
    - A flow steering moves a large amount of traffic
  - Less elements to deal with for load balancing but more difficult to balance

## Reactive vs. Proactive

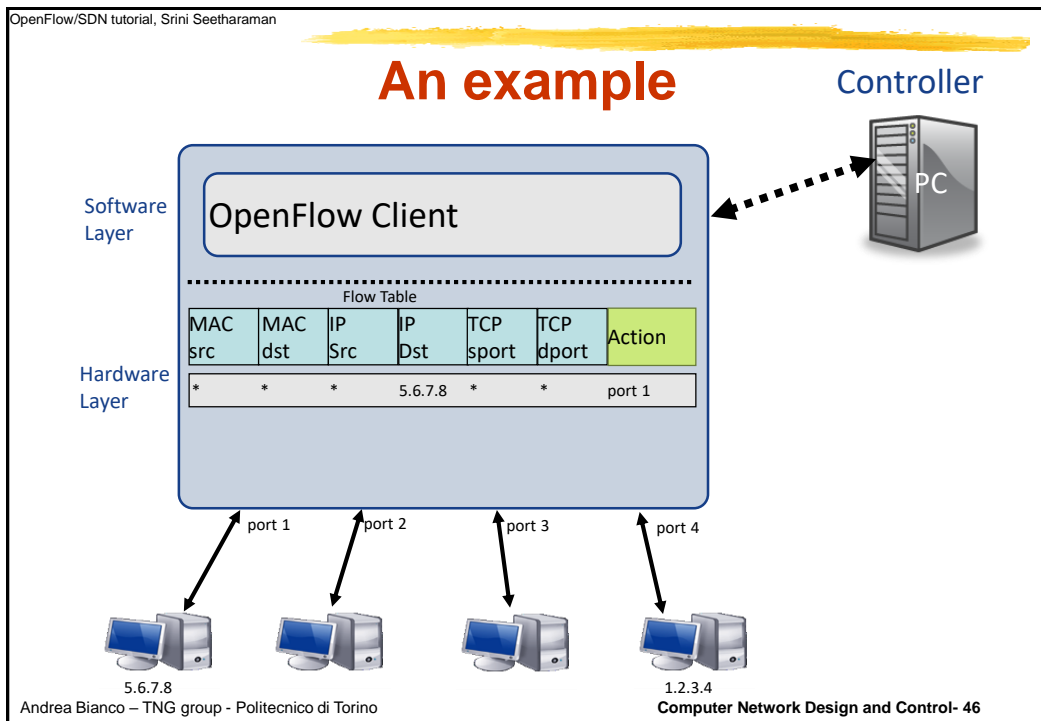
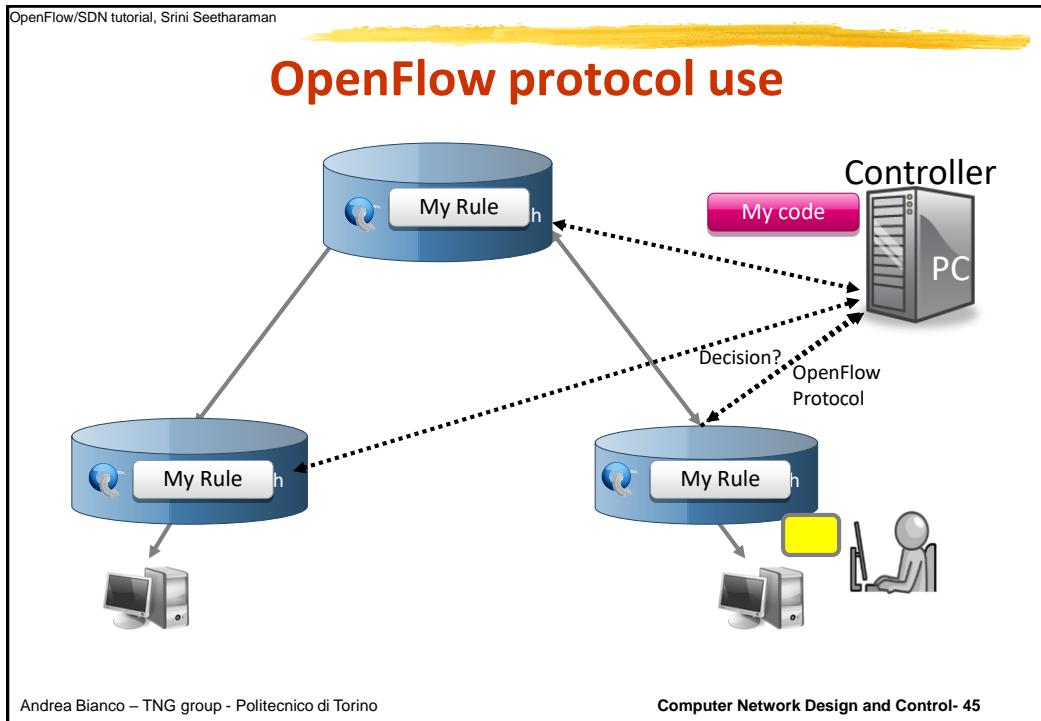
- Reactive
  - Flow table empty at boot
  - First packet of a flow sent to the controller
  - Controller inserts flow entries
  - Dynamic network
  
  - Every flow incurs small (?) additional flow setup time
  - Large control traffic
  - Large load on the controller
  - Efficient use of flow table
  - If control connection lost, switch has limited utility
- Proactive
  - Controller pre-populates flow table in switch at boot
  - Zero additional flow setup time
  - Static network
  
  - Loss of control connection does not disrupt traffic
  - Essentially requires aggregated (wildcard) rules
    - Reduced table size

## OpenFlow protocol

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# Software Defined Networking



## OpenFlow protocol messages

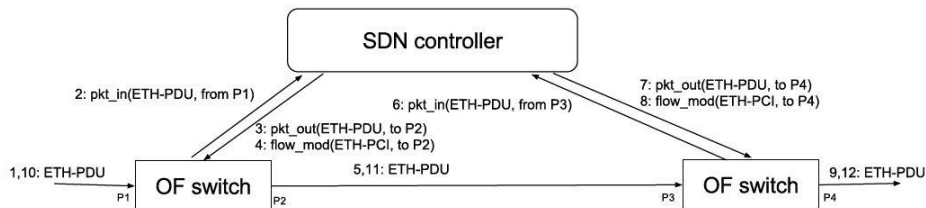
- **Controller-to-switch**
  - Initiated by the controller and used to directly manage or inspect the state of the switch
    - Features, Config, Modify State, Read State, Packet Out, Barrier
- **Asynchronous**
  - Sent to the controller without controller soliciting
    - Packet-in, Flow Removed/Expiration, Port status, Error, ...
- **Symmetric**
  - Sent without solicitation in any direction
    - Hello, Echo, Experimenter/Vendor

## OpenFlow (main) messages

- **Packet\_in**
  - Switch to controller
  - Carries a packet copy (possibly only the header)
    - What is best?
  - Generated by default in case of table miss
- **Packet\_out**
  - Controller to switch
  - Send the packet out of a specified port
  - Carries the full packet or the switch buffer id
- **Flow\_mod**
  - Controller to switch
  - Modify flow tables
  - Carries match-action rule to install



## OpenFlow example

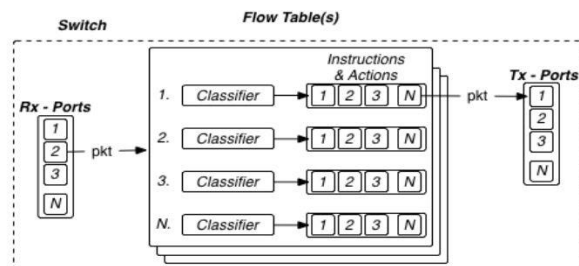


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## Packet processing

- Packets arrive and leave through ports
- Packets are matched to flow in flow tables using classifiers
- Flows contain set of instructions and actions applied to each packet in the match

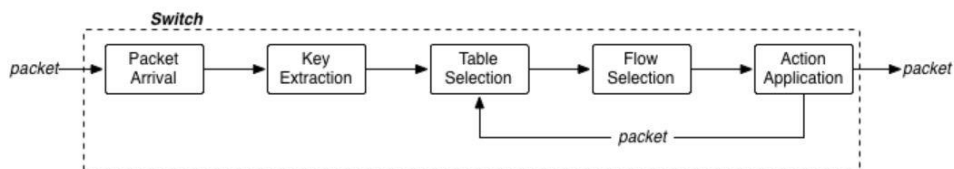


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## Packet lifecycle

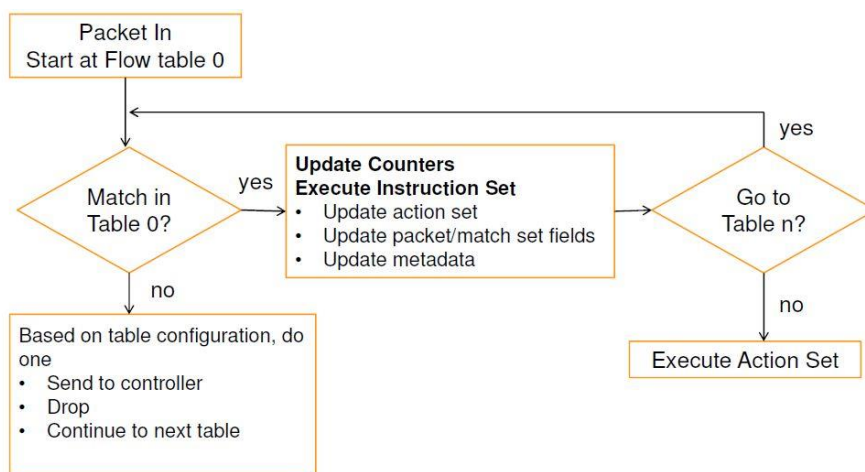
- On packet arrival a key is built
  - Metadata (arrival time, arrival port, memory location)
  - Fields in packet header
- Key is used to select a flow in the table
- Actions associated with the flow are applied
  - Drop, mutate, queue, forward, move to next table



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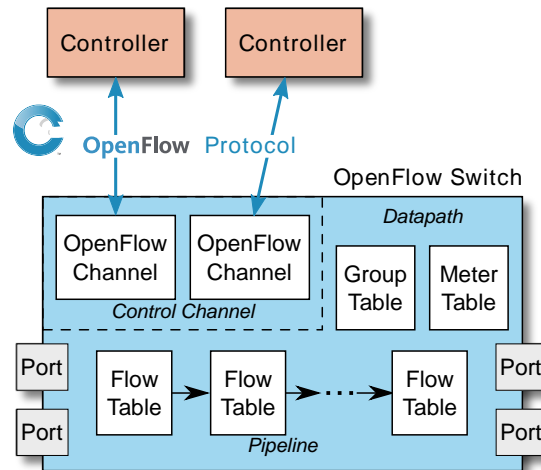
## Packet matching



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## Openflow switch implementation

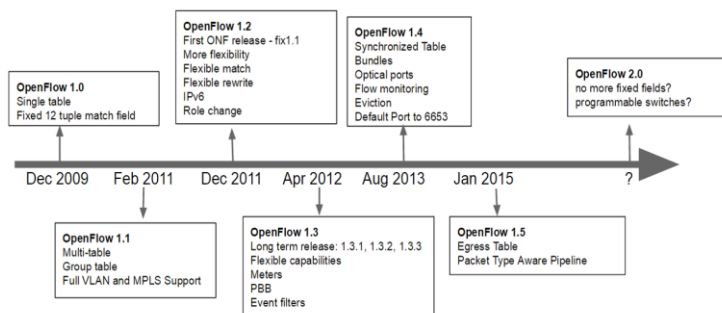


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## Openflow versions

- Published by Open Networking Foundation
  - No profit
  - Funded by Deutsche Telekom, Facebook, Google, Microsoft, Verizon, etc.



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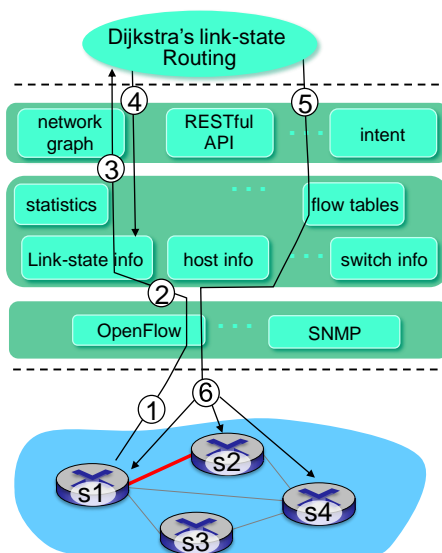
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## SDN architecture in action

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From Kurose Ross: Computer Networking

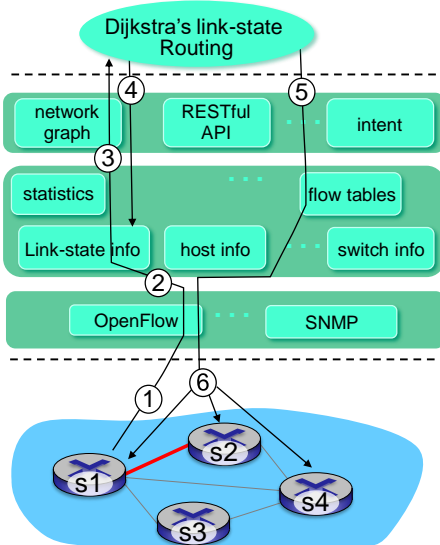
## An example



- ① S1, experiencing link failure using OpenFlow port status message to notify controller
- ② SDN controller receives OpenFlow message, updates link status info
- ③ Dijkstra's routing algorithm application has previously registered to be called when ever link status changes. It is called.
- ④ Dijkstra's routing algorithm access network graph info, link state info in controller, computes new routes

From Kurose Ross: Computer Networking

## An example



- ⑤ Link state routing app interacts with flow-table-computation component in SDN controller, which computes new flow tables needed
- ⑥ Controller uses OpenFlow to install new tables in switches that need updating

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## Distributed controllers

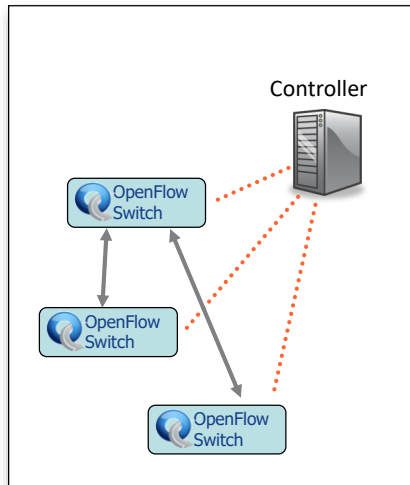
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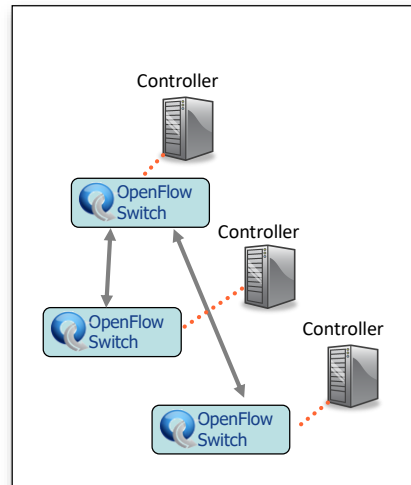
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## Centralized vs Distributed Control

### Centralized Control



### Distributed Control



## Why distributed/multiple controllers?

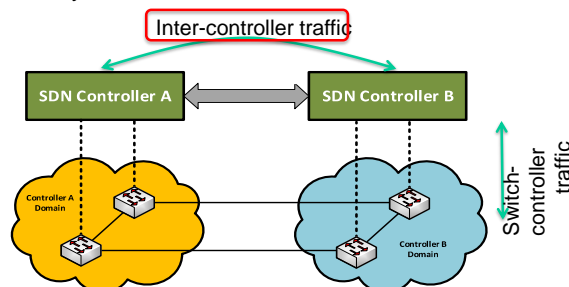
- To enhance resilience to failures
  - Controller failures can be managed
  - Still to deal with failures in data and control plane
- To solve scalability issues
  - Faster controllers
    - Limited scaling
  - More proactive rules to reduce number of requests
    - Limited flexibility
  - Multiple controllers
    - Permit load balancing to reduce processing load
    - Permit switch migration

## Distributed controllers

- Virtual topology among controllers
  - to coordinate the operations of the controllers
  - peer, hierarchical, master/slave
- Network view maintenance
  - different levels of consistency (strong/weak) among the controllers
  - affects the reactivity
  - may lead to temporary rule conflicts

## Control plane in distributed controllers

- Switch-controller (Sw-Ctr) traffic
  - Standardized
- Controller-controller (Ctr-Ctr) traffic (East-West-bound interfaces)
  - Proprietary
  - To get consistent view
  - May be non negligible
  - Critical for reactivity

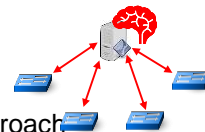


## Stateful data plane

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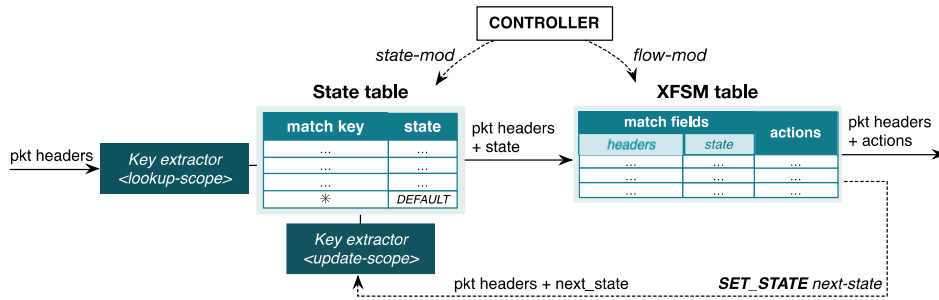
## Stateful SDN dataplane

- Stateless approach (OpenFlow)
  - Stateless switches, all the states in the controller
  - Limited reactivity due to the (logically) centralized approach
- Stateful approach: OpenState, OpenPacketProcessor (OPP), P4
  - Permit some level of stateful processing (e.g., finite state machines) within switches
    - OpenState adds a state table (IF state A THEN IF state B THEN)
    - OpenPacketProcessor: state defined with multiple variables, counters,
    - P4 much more flexible (description language of HW behavior)
  - Enabled by new generation of hardware
    - 6.5Tbps Tofino chipset @ Barefoot Networks





## Hardware implementation

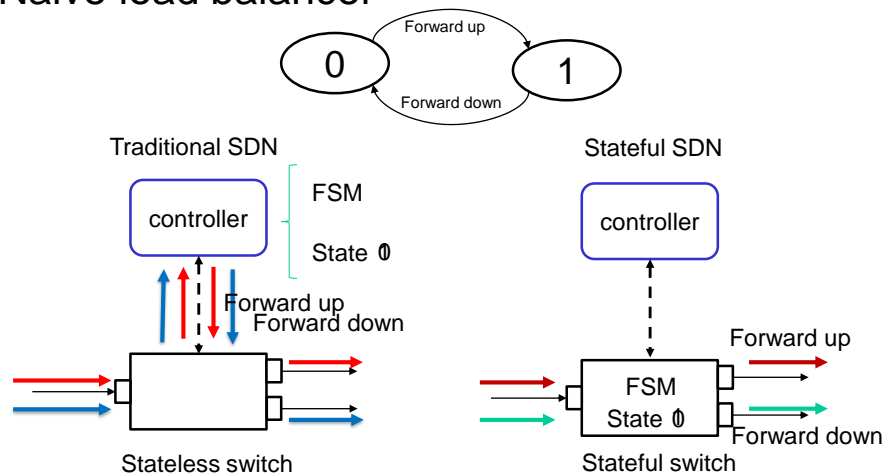


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## Toy example

- Naive load balancer

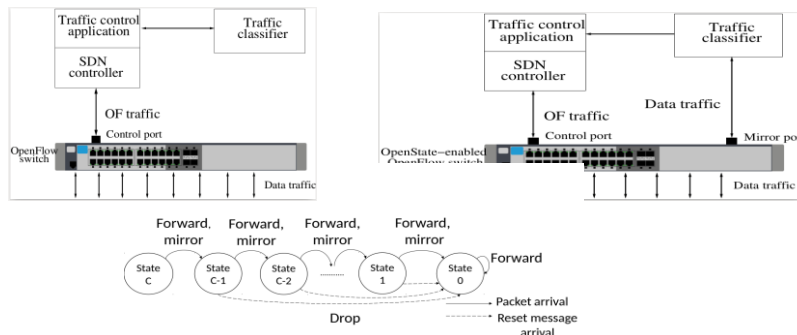


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## Traffic classification

- Mirror a pre-defined number of packets to traffic classifier for each flow
- Interrupt the mirroring if the flow is identified



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## Stateful benefits

- Improve network reactivity
  - Simple local decisions at the switch
  - Reduced controller load
  - Reduced signaling overhead
- Permits to gracefully move functionalities
  - Balance central vs distributed control
- Not all switches need to be stateful
  - State positioning or distribution

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Computer Network Design and Control- 68